

Signature schemes

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Introduction

A signature scheme is a method of signing a message stored in electronic form.

Conventional signature

Islam put an end to the idea of a generation, father to son, from

So please in Allah's name, do justice and mercy which Islam

Wa as Salamun 'Alaykum



Yusuf Islam

Digital signature

```
Signature: [c=]
12583171957620269831141560116
7093217610124910902417569020303
319933409810 [d=]
77901292515016200100658291
1504547126084771614879514744953
619459824421453
Signature length: 476
Algorithm: ECSP-DSA Alg
Hash function:
SHA-1
Key: [HybridEncryption][B
ob][EC-prime239v1][1178702474][
PIN=1234] Me
ssage: Cryptool....Cryptool
1 is a comprehensive educationa
1 program about cryptography an
d cryptanalysis.....This is a t
ext file, shown in order to hel
p you to make your first steps
with Cryptool.....1) As a first
```

The signature scheme

Definition 7.1: A *signature scheme* is a five-tuple $(\mathcal{P}, \mathcal{A}, \mathcal{K}, \mathcal{S}, \mathcal{V})$, where the following conditions are satisfied:

1. \mathcal{P} is a finite set of possible *messages*
2. \mathcal{A} is a finite set of possible *signatures*
3. \mathcal{K} , the *keyspace*, is a finite set of possible *keys*
4. For each $K \in \mathcal{K}$, there is a *signing algorithm* $\mathbf{sig}_K \in \mathcal{S}$ and a corresponding *verification algorithm* $\mathbf{ver}_K \in \mathcal{V}$. Each $\mathbf{sig}_K : \mathcal{P} \rightarrow \mathcal{A}$ and $\mathbf{ver}_K : \mathcal{P} \times \mathcal{A} \rightarrow \{true, false\}$ are functions such that the following equation is satisfied for every message $x \in \mathcal{P}$ and for every signature $y \in \mathcal{A}$:

$$\mathbf{ver}_K(x, y) = \begin{cases} true & \text{if } y = \mathbf{sig}_K(x) \\ false & \text{if } y \neq \mathbf{sig}_K(x). \end{cases}$$

A pair (x, y) with $x \in \mathcal{P}$ and $y \in \mathcal{A}$ is called a *signed message*.

Example

Cryptosystem 7.1: RSA Signature Scheme

Let $n = pq$, where p and q are primes. Let $\mathcal{P} = \mathcal{A} = \mathbb{Z}_n$, and define

$$\mathcal{K} = \{(n, p, q, a, b) : n = pq, p, q \text{ prime}, ab \equiv 1 \pmod{\phi(n)}\}.$$

The values n and b are the public key, and the values p, q, a are the private key.

For $K = (n, p, q, a, b)$, define

$$\mathbf{sig}_K(x) = x^a \pmod{n}$$

and

$$\mathbf{ver}_K(x, y) = \text{true} \Leftrightarrow x \equiv y^b \pmod{n}$$

$(x, y \in \mathbb{Z}_n)$.

Types and goals of attacks

Types of attacks

- Key-only
- Known message
- Chosen message

Goals of attacks

- Total break
- Selective forgery
- Existential forgery

An example of existential attack

Ex: Textbook RSA signature

$$\mathbf{sig}_K(x) = x^a \bmod n \quad \mathbf{ver}_K(x, y) = \text{true} \Leftrightarrow x \equiv y^b \pmod{n}$$

C1: Choose any signature y , then compute $x = y^b \bmod n$.

C2:

$y_1 = \mathbf{sig}_K(x_1)$ and $y_2 = \mathbf{sig}_K(x_2)$ are any two messages previously signed

$$\mathbf{ver}_K(x_1 x_2 \bmod n, y_1 y_2 \bmod n) = \text{true},$$

→ Choose $y_2 = y/y_1$.

→ Need to be hashed.